# Discovery through Automation

Luca Cardelli, University of Oxford QAVAS Oxford, 2021-09-28

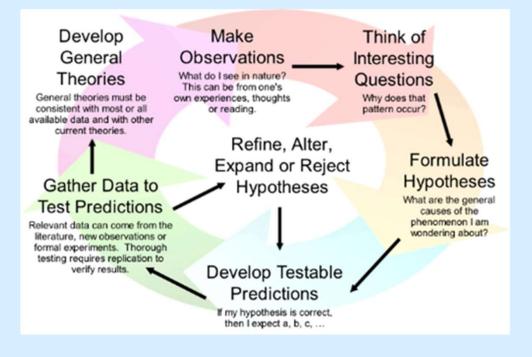
# Outline

- The Scientific Method Its eventual automation
- Models (that know nothing about protocols)
  - **Chemical Reaction Networks**
- Lab Protocols (that know nothing about models)
   Digital Microfluidics
- Integration
  - Closed-loop modeling and protocol execution The Kaemika App

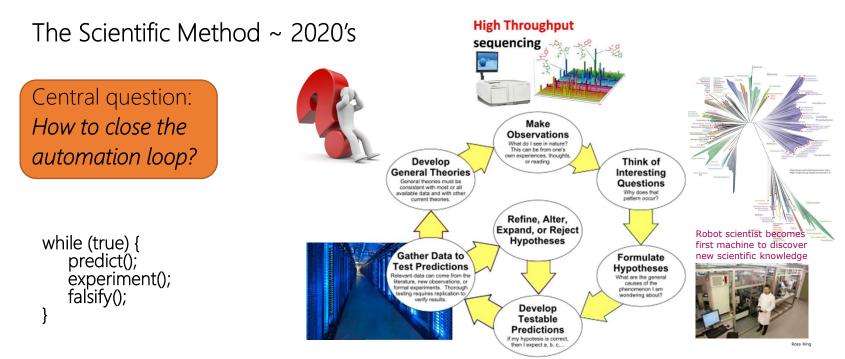
### The Scientific Method

Hasan Ibn al-Haytham (1027) Book of Optics

Galileo Galilei (1638) Two New Sciences



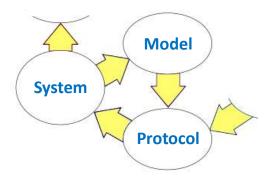
#### Discovery through Automation



Garland, Jr., Theodore. "The Scientific Method as an Ongoing Process". U C Riverside.

#### The Inner Loop

- A model is refined by testing a (fixed) protocol against a systems
- A *protocol* is refined by testing a (fixed) *model* against a *systems*
- Today: publication does not accurately reflect execution
  - Model: poorly-maintained matlab script
  - Protocol: poorly-described manual steps in the lab
  - System: poorly-characterized and hardly "resettable"
  - ⇒ Crisis in biology: experiments are done once and are hard to reproduce http://www.nature.com/news/reproducibility-1.17552



#### Models

(those things that know nothing about protocols)

#### Chemical Reaction Networks (CRN)

 $X + Y \rightarrow r Z + W$ 

- A *phenomenological model* of kinetics in the natural sciences By (only) observing naturally occurring reactions
- A *programming language*, *finitely* encoded in the genome By which living things manage the *unbounded* processing of matter and information
- A *mathematical structure*, rediscovered in many forms Vector Addition Systems, Petri Nets, Bounded Context-Free Languages, Population Protocols, ...
- A description of *mechanism* ("instructions" / "interactions") rather than *behavior* ("equations" / "approximations")
  - Although the two are related in precise ways
  - Enabling, e.g., the study of the evolution of mechanism through unchanging behavior

#### Programming Examples spec program Y := 2X $X \rightarrow Y + Y$ $Y := \lfloor X/2 \rfloor$ $X + X \rightarrow Y$ Y := X1 + X2X1 -> Y X2 -> Y Y := min(X1, X2) X1 + X2 -> Y

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Advanced Programming Examples		
spec	program	
Y := max(X1, X2)	X1 -> L1 + Y X2 -> L2 + Y	max(X1,X2)= (X1+X2)-min(X1,X2)
	L1 + L2 -> K Y + K -> 0	(but is not computed "sequentially")
Approximate Majority		
(X,Y) := if X≥Y then (X+Y, 0) if Y≥X then (0, X+Y)	X + Y -> Y + B Y + X -> X + B B + X -> X + X B + Y -> Y + Y	

Programming<sup>v</sup>any algorithm as a CRN

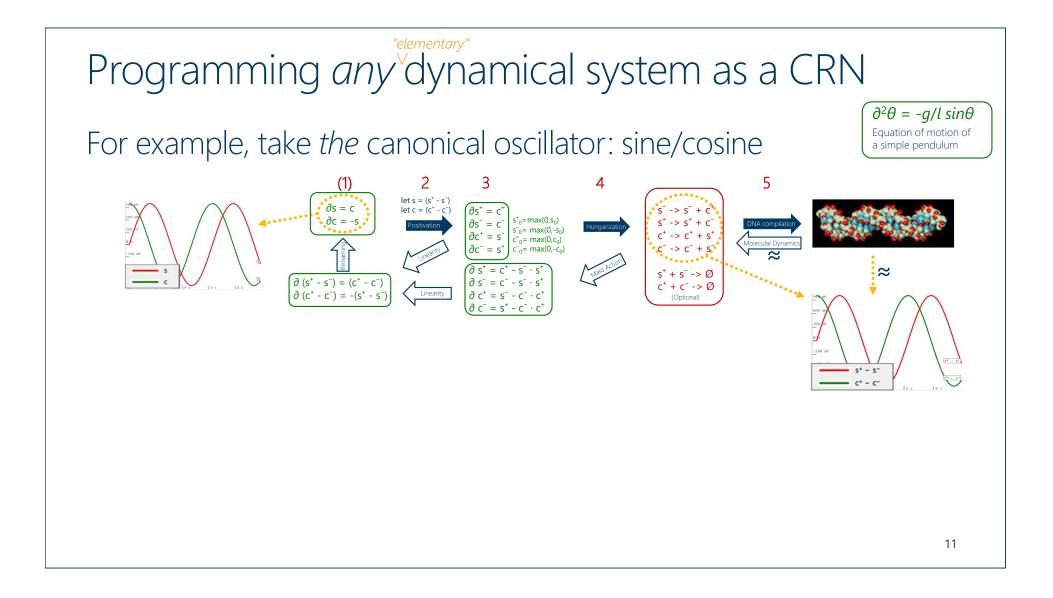
A CRN is a *finite* set of reactions over a *finite* set of species

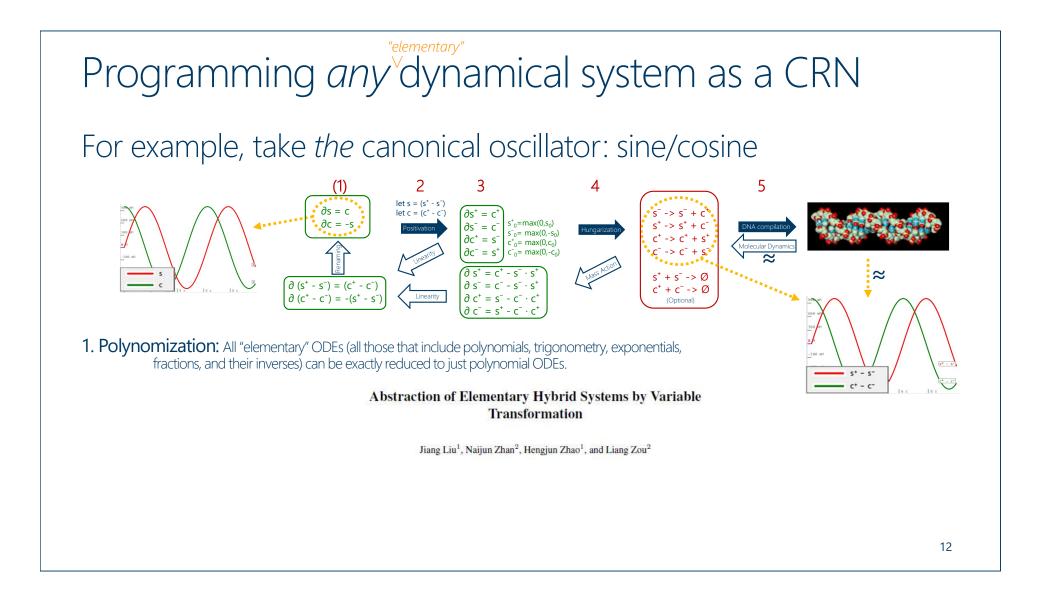
CRNs are not Turing complete Like Petri nets: reachability is decidable

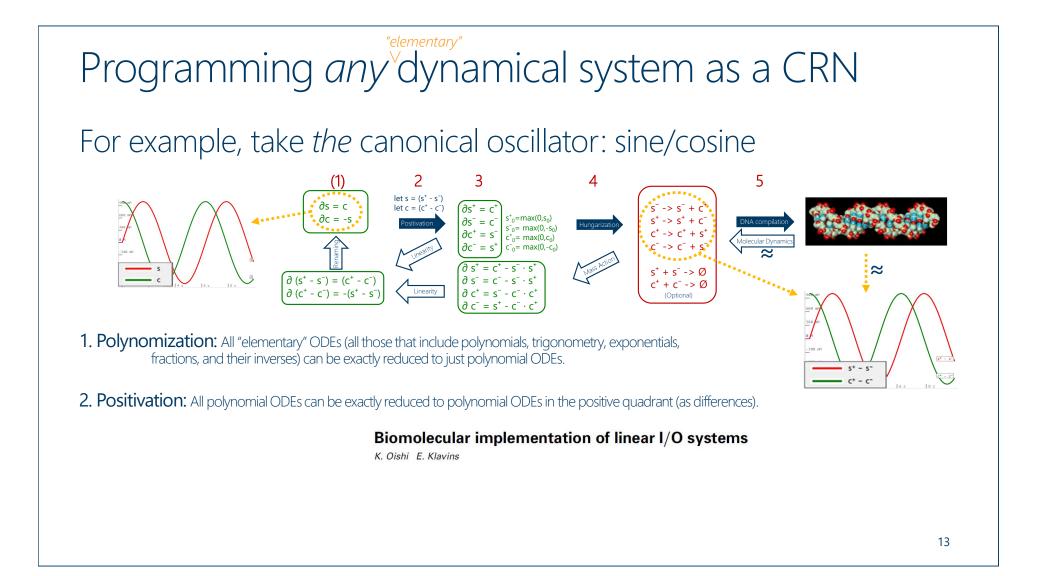
#### But unlike Petri nets, CRNs are *approximately* Turing complete Because reactions have also *rates*

This make it possible to approximate Turing completeness by approximating test-for-zero in a register machine. The probability of error (in test-for-zero) can be made arbitrarily small over the entire (undecidably long) computation.

Adding polymerization to the model makes it fully Turing complete

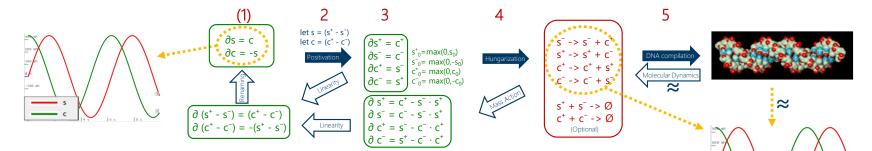








#### For example, take the canonical oscillator: sine/cosine

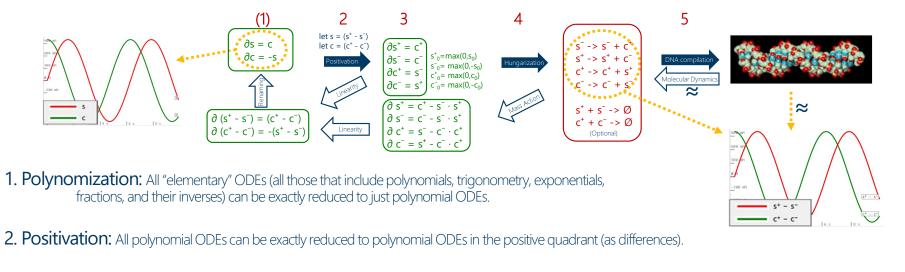


- 1. Polynomization: All "elementary" ODEs (all those that include polynomials, trigonometry, exponentials, fractions, and their inverses) can be exactly reduced to just polynomial ODEs.
- 2. Positivation: All polynomial ODEs can be exactly reduced to polynomial ODEs in the positive quadrant (as differences).
- 3. All positivized ODEs are Hungarian: I.e., all negative monomials have their I.h.s. differential variable as a factor.

s+ - sc+ - c-



#### For example, take the canonical oscillator: sine/cosine



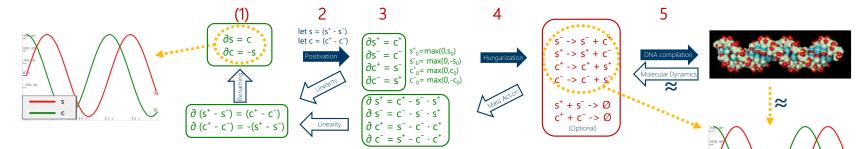
- 3. All positivized ODEs are Hungarian: I.e., all negative monomials have their I.h.s. differential variable as a factor.
- 4. Hungarization: All Hungarian ODEs can be exactly reduced to mass action CRNs.

ON THE INVERSE PROBLEM OF REACTION KINETICS

V. HÁRS - J. TÓTH



#### For example, take the canonical oscillator: sine/cosine



- 1. Polynomization: All "elementary" ODEs (all those that include polynomials, trigonometry, exponentials, fractions, and their inverses) can be exactly reduced to just polynomial ODEs.
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- 4. Hungarization: All Hungarian ODEs can be exactly reduced to mass action CRNs.

#### DNA as a universal substrate for chemical kinetics

s+ - sc+ - c-

David Soloveichik, Georg Seelig, and Erik Winfree PNA5 March 23, 2010 107 (12) 5393-5398; https://doi.org/10.1073/pnas.090938010

5. Molecular Programming: All mass action CRNs, up to time rescaling, can be arbitrarily approximated by engineered DNA molecules.

Chemistry is (also) a formal language that we can use to implement *any* algorithm and *any* dynamical system with *real* (DNA) molecules

Turing complete and "Shannon complete"

- ANY collection of abstract chemical reactions can be implemented with specially designed DNA molecules, with accurate kinetics (up to time scaling).
- Approaching a situation where we can "systematically compile" (synthesize) a model to DNA molecules, run an (automated) protocol, and observe (sequence) the results in a closed loop.

### A Model

A Chemical Reaction Network, provided explicitly or (in this case) generated from a higher-level description of the initial strands, according to the DNA strand displacement rules

$$\frac{2}{1*} \frac{3}{2*} \frac{4*}{3*} + \frac{2}{1} \frac{3}{2} \frac{60.0003}{1*} \frac{2}{2*} \frac{3}{3*} \frac{4}{4*}$$

$$\frac{2}{1*} \frac{3}{2*} \frac{4}{3*} + \frac{1}{12} \frac{2}{2} \frac{0.0003}{2} \frac{1}{1*} \frac{2}{2*} \frac{3}{3*} \frac{4}{4*}$$

$$\frac{1}{1*} \frac{2}{2*} \frac{3}{3*} \frac{4}{4*} + \frac{3}{4} \frac{0.0003}{2} \frac{1}{1*} \frac{2}{2*} \frac{3}{3*} \frac{4}{4*} + \frac{2}{2} \frac{3}{4}$$

$$\frac{2}{1*} \frac{3}{2*} \frac{4}{3*} + \frac{3}{4} \frac{0.0003}{2} \frac{1}{1*} \frac{2}{2*} \frac{3}{3*} \frac{4}{4*} + \frac{2}{2} \frac{3}{4}$$

$$\frac{2}{1*} \frac{3}{2*} \frac{4}{3*} + \frac{3}{4*} \frac{4}{2} \frac{0.0003}{2} \frac{1}{1*} \frac{2}{2*} \frac{3}{3*} \frac{4}{4*} + \frac{2}{2} \frac{3}{4}$$

# Model Semantics (deterministic)

#### ODE semantics of CRNs

State produced by a CRN C = (A, R) (species A, reactions R) with flux F (r.h.s. of its mass action ODEs) at time t, from initial state  $(x_0, V, T)$  (initial concentrations  $x_0$ , volume V, temperature T):

 $\llbracket ((\mathcal{A},\mathcal{R},x_0),V,T) \rrbracket(H)(t) = (G(t),V,T)$ let  $G : [0...H) \to \mathbb{R}^{|\mathcal{A}|}$  be the solution of  $G(t') = x_0 + \int_0^{t'} F(V,T)(G(s)) ds$ 

#### Summarizing

- Our models are (chemical) programs
- We can compute their behavior (their final state)
- $\cdot\,$  We can (virtually) run them by integration of the ODEs
- $\cdot$  We can (physically) run them by DNA nanotech

#### Protocols

(those things that know nothing about models)

#### A Protocol For DNA gate assembly and activation in vitro 2 3 Sample ()2\* 3\* Gate<sub>B</sub> 1\* 4\* Output Mix Protocol steps Dispense (liquid handing) Equilibrate 2 Dispose Input<sub>1</sub> $p_1$ $p_2$ t<sub>1</sub> 3 Input<sub>2</sub> 3 2 p<sub>3</sub> 1\* 2\* 3\* 4\* Gate Ř₄ t<sub>2</sub>

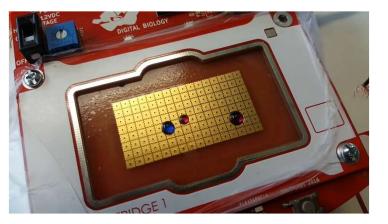
# Digital Microfluidics

#### OpenDrop

https://www.youtube.com/watch?v=ncfZWqPm7-4



OpenDrop speed test https://www.youtube.com/watch?v=pSls9L h3Q0



Purple Drop (UW) https://misl.cs.washington.edu/projects/fluidics.html



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# Digital Microfluidics

- A general, *programmable*, platform to execute the main liquid-handling operations
- To close the cycle, it can support many automated observation techniques on-board or off-board via peripheral pumps (sequencing, mass spec, ...) although these are all very hardware-dependent.

### A Protocol Language

Samples: containers with volume, temperature, concentrations

P =

 $\begin{array}{lll} x & (a \; sample \; variable) \\ & (x_0,V,T) & (initial \; condition) \\ & let \; x = P_1 \; in \; P_2 & (define \; local \; variable) \\ & Mix(P_1,P_2) & (mix \; samples) \\ & let \; x,y = Split(P_1,p) \; in \; P_2 & (split \; samples) \\ & Equilibrate(P,t) & (equilibrate \; sample \; for \; t \; seconds) \\ & Dispose(P) & (discard \; sample) \end{array}$ 

Experimental Biological Protocols with Formal Semantics

Alessandro Abate², Luca Cardelli<sup>1,2</sup>, Marta Kwiatkowska², Luca Laurenti², and Boyan Yordanov<sup>1</sup>

 $^{1}\,$  Microsoft Research Cambridge  $^{2}\,$  Department of Computer Science, University of Oxford

#### Protocol Semantics (deterministic)

Each program denotes a *final* state < concentrations, volume, temperature>

 $\llbracket P \rrbracket^{
ho}$  is the final state produced by a protocol P where ho binds its free variables:

$$\begin{split} \|x\|^{\rho} &= \rho(x) \\ \|x_{0}, V, T\|^{\rho} &= (x_{0}, V, T) \\ \|Mix(P_{1}, P_{2})\|^{\rho} &= \\ let(x_{0}^{1}, V_{1}, T_{1}) &= \|P_{1}\|^{\rho} \\ let(x_{0}^{2}, V_{2}, T_{2}) &= \|P_{2}\|^{\rho} \\ let(x_{0}^{2}, V_{2}, T_{2}) &= \|P_{2}\|^{\rho} \\ (\frac{x_{0}^{1}V_{1} + x_{0}^{2}V_{2}}{V_{1} + V_{2}}, V_{1} + V_{2}, \frac{T_{1}V_{1} + T_{2}V_{2}}{V_{1} + V_{2}}) \\ \|letx &= P_{1} in P_{2}\|^{\rho} &= \\ let(x_{0}, V, T) &= \|P_{1}\|^{\rho} \\ let(x_{0$$

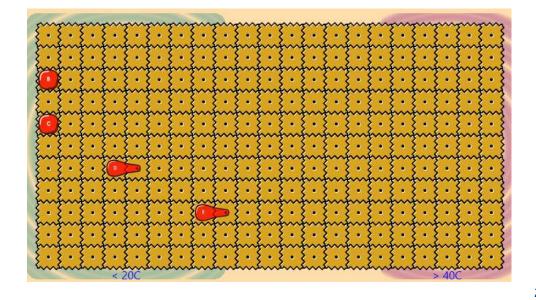
### Kaemika Microfluidics Compiler

- Mix, split, equilibrate, dispose
- Automatic routing no geometrical information
- Hot/cold zones

sample A {3µL, 20C}

split B,C,D,E = A
mix F = E,C,B,D

dispose F



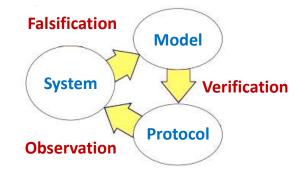
#### Summarizing

- Our protocols are (liquid handling) programs
- We can compute their behavior (their final state)
- We can (virtually) run them (by simulation)
- We can (physically) run them (by digital microfluidics)

### Models together with Protocols

## Automating "the whole thing"

- Protocols: sets of steps to direct lab machinery (or people)
  - Published in specialized journals. With varying accuracy.
- · Models: sets of equations to predict the results of lab experiments
  - Published in Auxiliary Online Materials. With lots of typos.
- Protocols know nothing about models
  - What hypothesis is the protocol trying to test? It is not written in the protocol.
- Models know nothing about protocols
  - What lab conditions are being used to test the model? It is not written in the model.
- While presumably talking about the same system
  - Through the experiment.
- Reproducibility crisis
  - Experiments are hard to reproduce. (materials, conditions, shortcuts)
  - · Even models are hard to reproduce! (typos in equations, sketchy diagrams, unexplained graphs, mysterious scripts)
- Similar to classical lifecycle problems in C.S.
  - · Documentation (model) gets out of step from code (protocol) if their integration is not automated.



# An Integrated Description

Samples: containers with volume, temperature, concentrations

P =

 $\begin{array}{lll} x & (a \; sample \; variable) \\ & (x_0,V,T) & (initial \; condition) \\ & let \; x = P_1 \; in \; P_2 & (define \; local \; variable) \\ & Mix(P_1,P_2) & (mix \; samples) \\ & let \; x,y = Split(P_1,p) \; in \; P_2 & (split \; samples) \\ & Equilibrate(P,t) & (equilibrate \; sample \; for \; t \; seconds) \\ & Dispose(P) & (discard \; sample) \end{array}$ 

each sample evolves (via *Equilibrate*) according to a given overall CRN:

 $\mathcal{C} = (\mathcal{A}, \mathcal{R})$  (species, reactions)

#### Experimental Biological Protocols with Formal Semantics

Alessandro Abate², Luca Cardelli<sup>1,2</sup>, Marta Kwiatkowska², Luca Laurenti², and Boyan Yordanov<sup>1</sup>

<sup>1</sup> Microsoft Research Cambridge
<sup>2</sup> Department of Computer Science, University of Oxford



#### Joint script

 $Input_1 = <1^* 2 > Output = <2 3 > Input_2 = <3 4^* > Gate = \{1^*\}[2 3]\{4^*\}$ 

$$\begin{split} P_1 =& let \, In1 = ((Input1, 100.0nM), 0.1mL, 25.0^\circ C) \, in \\ let \, In2 = ((Input2, 100.0nM), 0.1mL, 25.0^\circ C) \, in \\ let \, GA = ((Output, 100.0nM), 0.1mL, 25.0^\circ C) \, in \\ let \, GB = ((Gate_B, 100.0nM), 0.1mL, 25.0^\circ C) \, in \\ let \, sGA_{,=} \, Dispense(GA, p_1) \, in \\ let \, sGB_{,=} \, Dispense(GB, p_2) \, in \\ let \, sIn1_{,=} \, Dispense(In1, p_3) \, in \\ let \, sIn2_{,=} \, Dispense(In1, p_4) \, in \\ Observe(Equilibrate(Mix(Mix(Equilibrate(Mix(sGA, sGB), t_1), sIn2), t_2), idn). \end{split}$$

### Program Semantics (deterministic)

Each program denotes a *final* state < concentrations, volume, temperature>

 $[\![P]\!]^{
ho}$  is the final state produced by a protocol  $\ P$  for a fixed CRN  $\ {\cal C} = ({\cal A}, {\cal R})$  :

$$\begin{split} \llbracket x \rrbracket^{\rho} &= \rho(x) \\ \llbracket x_{0}, V, T \rrbracket^{\rho} &= (x_{0}, V, T) \\ \llbracket Mix(P_{1}, P_{2}) \rrbracket^{\rho} &= \\ let(x_{0}^{1}, V_{1}, T_{1}) &= \llbracket P_{1} \rrbracket^{\rho} \\ let(x_{0}^{2}, V_{2}, T_{2}) &= \llbracket P_{2} \rrbracket^{\rho} \\ let(x_{0}^{2}, V_{2}, T_{2}) &= \llbracket P_{2} \rrbracket^{\rho} \\ (\frac{x_{0}^{1}V_{1} + x_{0}^{2}V_{2}}{V_{1} + V_{2}}, V_{1} + V_{2}, \frac{T_{1}V_{1} + T_{2}V_{2}}{V_{1} + V_{2}}) \\ \llbracket let(x_{0}, V, T) &= \llbracket P_{1} \rrbracket^{\rho} \\ [letx = P_{1} in P_{2} \rrbracket^{\rho} = \\ let(x_{0}, V, T) &= \llbracket P_{1} \rrbracket^{\rho} \\ \llbracket let(x_{0}, V, T) &= \llbracket P_{1} \rrbracket^{\rho} \\ [letx = P_{1} in P_{2} \rrbracket^{\rho} = \\ let(x_{0}, V, T) &= \llbracket P_{1} \rrbracket^{\rho} \\ [letx = P_{1} in P_{2} \rrbracket^{\rho} = \\ let(x_{0}, V, T) &= \llbracket P_{1} \rrbracket^{\rho} \\ [letx = P_{1} in P_{2} \rrbracket^{\rho} = \\ let(x_{0}, V, T) &= \llbracket P_{1} \rrbracket^{\rho} \\ [letx = P_{1} in P_{2} \rrbracket^{\rho} = \\ let(x_{0}, V, T) &= \llbracket P_{1} \rrbracket^{\rho} \\ [letx = P_{1} in P_{2} \rrbracket^{\rho} = \\ let(x_{0}, V, T) &= \llbracket P_{1} \rrbracket^{\rho} \\ [letx = P_{1} in P_{2} \rrbracket^{\rho} = \\ let(x_{0}, V, T) &= \llbracket P_{1} \rrbracket^{\rho} \\ [letx = P_{1} in P_{2} \rrbracket^{\rho} = \\ let(x_{0}, V, T) &= \llbracket P_{1} \rrbracket^{\rho} \\ [letx = P_{1} in P_{2} \rrbracket^{\rho} = \\ let(x_{0}, V, T) &= \llbracket P_{1} \rrbracket^{\rho} \\ [letx = P_{1} in P_{2} \rrbracket^{\rho} = \\ let(x_{0}, V, T) &= \llbracket P_{1} \rrbracket^{\rho} \\ [letx = P_{1} in P_{2} \rrbracket^{\rho} = \\ let(x_{0}, V, T) &= \llbracket P_{1} \rrbracket^{\rho} \\ [letx = P_{1} in P_{2} \rrbracket^{\rho} = \\ let(x_{0}, V, T) &= \llbracket P_{1} \rrbracket^{\rho} \\ [letx = P_{1} in P_{2} \rrbracket^{\rho} = \\ let(x_{0}, V, T) &= \llbracket P_{1} \rrbracket^{\rho} \\ [letx = P_{1} in P_{2} \rrbracket^{\rho} = \\ let(x_{0}, V, T) &= \llbracket P_{1} \rrbracket^{\rho} \\ [letx = P_{1} in P_{2} \rrbracket^{\rho} = \\ let(x_{0}, V, T) &= \llbracket P_{1} \rrbracket^{\rho} \\ [letx = P_{1} in P_{2} \rrbracket^{\rho} = \\ let(x_{0}, V, T) &= \llbracket P_{1} \rrbracket^{\rho} \\ [letx = P_{1} in P_{2} \rrbracket^{\rho} = \\ let(x_{0}, V, T) &= \llbracket P_{1} \rrbracket^{\rho} \\ [letx = P_{1} in P_{2} \rrbracket^{\rho} = \\ let(x_{0}, V, T) &= \llbracket P_{1} \rrbracket^{\rho} \\ [letx = P_{1} in P_{2} \rrbracket^{\rho} = \\ let(x_{0}, V, T) &= \llbracket P_{1} \rrbracket^{\rho} \\ [letx = P_{1} in P_{2} \rrbracket^{\rho} = \\ let(x_{0}, V, T) &= \llbracket P_{1} \rrbracket^{\rho} \\ [letx = P_{1} in P_{2} \rrbracket^{\rho} = \\ let(x_{0}, V, T) &= \llbracket P_{1} In P_{2} \rrbracket^{\rho} \\ [letx = P_{1} in P_{2} \rrbracket^{\rho} = \\ let(x_{0}, V, T) &= \llbracket$$

### A Joint Semantics

This semantics gives us a *joint simulation algorithm*, connecting chemical simulation with protocol simulation.

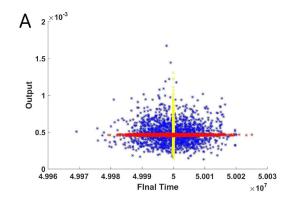
In this presentation everything is *deterministic*. The state of the protocol is passed to the chemical simulator, which computes a new state that it passes to the protocol simulator, and so on.

Kaemika uses such a joint simulation algorithm for *stochastic* simulation, passing also variance information back and forth between chemical and protocol simulation.

This requires an extension of the above semantics using the Linear Noise Approximation of chemical kinetics, which computes mean and variance of concentrations (both by ODEs, not e.g. by Gillespie algorithm), and a similar extension of the protocol operations.

### Stochastic Analysis

- We can ask: what is the probability of a certain outcome given uncertainties in *both the protocol and the model*?
- Conversely: which parameters of *both the protocol and the model* best fit the observed result?
- Also, we can use Statistical Modelchecking:



1500 executions including protocol uncertainty due timing and pipetting errors (red).

1500 executions including only model uncertainty about rates of the CRN (yellow).

1500 executions including both sources of uncertainty (blue).

We may estimate by Statistic Model Checking, e.g. the probability that Output will fall in a certain range, given distributions over uncertain model and protocol parameters. 34

# Simulating Reaction Networks together with Digital Protocols



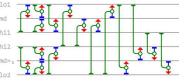


Search "Kaemika" in the app stores http://lucacardelli.name/kaemika.html An integrated language for chemical models & experimental protocols

Deterministic (ODE) and stochastic (LNA) simulation

Chemical reaction networks (CRNs) and liquid-handling protocols

Reaction scores



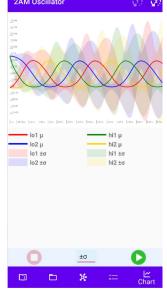
Functional scripting

GUI

# Kaemika

- A prototype language for chemical models & protocols
- <u>http://lucacardelli.name/kaemika.html</u>
- Search "Kaemika" in the App stores





- CRN simulation
- Microfluidics simulation
- Reaction graphs
- ODE equations
- Stochastic noise (LNA)

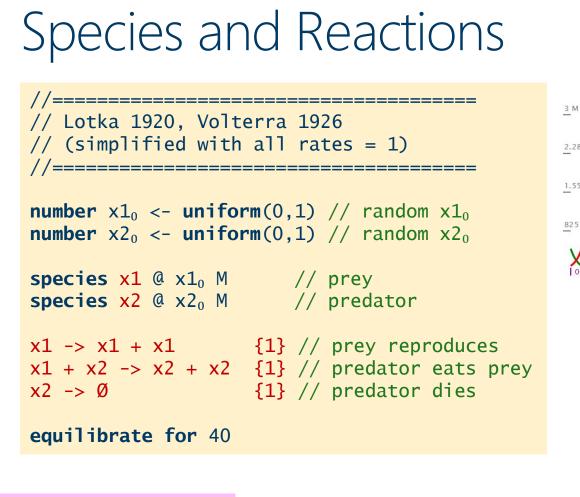
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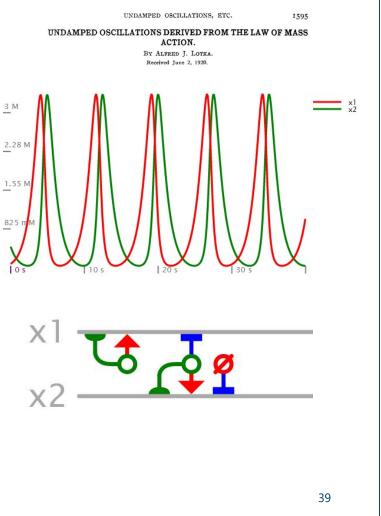
# Main features

- Species and reactions
  - · Characterized by initial values and rates
- "Samples" (compartments) and Protocols
  - $\cdot$  Isolate species and reactions in a compartment, and mix compartments

### • Kinetics (simulation)

- Deterministic (ODE) or stochastic (LNA) for chemical models
- Digital microfluidics for chemical protocols
- Programming abstractions
  - Assemble models and protocols as compositions of modules





<= Demo: LotkaVolterra

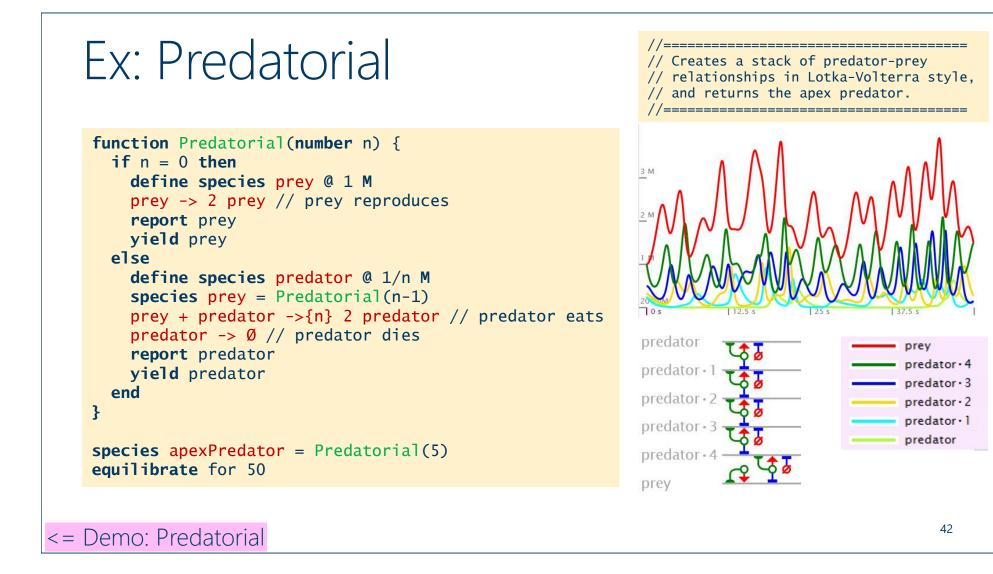
Stochastic (LNA) simulation
<ul> <li>For all programs (any CRN, any Protocol)</li> </ul>
$ \begin{array}{l} \\ \begin{array}{l} \\ \\ \\ \\ \\ \\ \\ \\ \\ \\ \\ \\ \\ \\ \\ \\ \\ \\ \\$

# Writing Models Compositionally

- Embedded chemical notation
  - Programs freely contain both chemical reactions and control flow Can generate unbounded-size reaction networks
- Rich data types
  - numbers, species, functions, networks, lists, flows (time-courses) flows are composable functions of time used in rates, plotting, and observation

### Modern abstractions

*Functional:* programs take *data* as parameters and produce *data* as results *Monadic:* programs also produce *effects* (*species, reactions, liquid handling*) *Nominal: lexically scoped* chemical species (species are not "strings")



# Describing a Protocol

- Samples (e.g., test tubes)
  - $\cdot\,$  Are characterized by a volume and a temperature
  - $\cdot\,$  Contain a specified set of species
  - $\cdot\,$  Evolve according to reactions that operates on those species
  - $\boldsymbol{\cdot}$  Isolate species and reactions
- Protocol Operations (e.g., liquid handling)
  - $\cdot\,$  Accept and produce samples
  - Accepted samples are used up (they can only be operated-on once)

# Samples

- · Samples contain concentrations of species, acted over by reactions.
- Each sample has a fixed volume and a fixed temperature through its evolution.
- Sample concentrations are in units of molarity M = mol/L.
- The default implicit sample is called the vessel {1 mL, 20 C}

```
species {c} // a species for multiple samples
```

```
sample A {1\muL, 20C} // volume and temperature
species a @ 10mM in A // species local to A
amount c @ 1mM in A // amount of c in A
a + c -> a + a
```

```
sample B {1\muL, 20C}
species b @ 10mM in B // species local to B
amount c @ 1mM in B // amount of c in B
b + c -> c + c
```

An amount can also be given in grams (if molar mass is specified). The resulting concentration is then <u>relative to sample volume</u>.

species {NaCl#58.44}

```
sample C {1mL, 20C}
amount NaCl @ 8g in C
```

```
Reactions can be specified with
Arrhenius parameters {collision
frequency, activation energy}.
The reaction kinetics is then
relative to sample temperature T.
```

a + c ->{2, 5} a + a // rate is 2\*e^(-5/(R\*T))

<= Demo: MixAndSplit

### Ex: Serial Dilution (recursive protocol)

```
network SerialDilution(number count, sample s, network f) {
    if count > 0 then
        sample solvent {9*observe(volume,s) L, observe(kelvin,s) K}
    mix s = s, solvent
    split s, dilution = s by 0.1, 0.9
    f(dilution)
    SerialDilution(count-1, s, f)
    end
}
```

//initial sample to be diluted:

```
sample init {1mL, 25C}
species A @ 1M in init
species B @ 1M in init
A + B ->{20} A
A -> Ø
```

```
//apply this network to each dilution;
//note that this invokes a simulation
//each time in each solution
```

```
network test(sample s) {
    equilibrate s for 10
    dispose s
}
```

```
//dilute 4 times
```

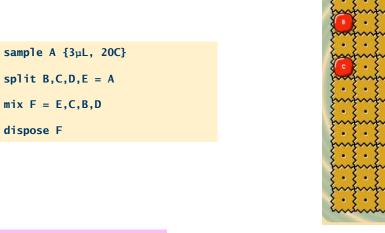
SerialDilution(4, init, test)

Prepare a series of increasingly diluted solutions and apply a network f to each (f can add species and reactions to the solutions)

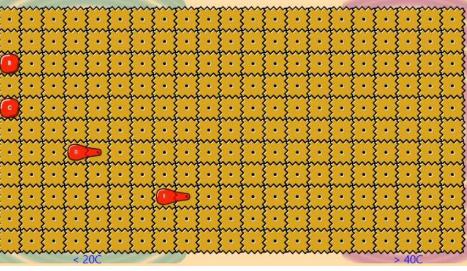
```
RESULT:
sample init {1mL, 298.2K} {A = 1M, B = 1M}
sample s2 {1mL, 298.2K} {A = 100mM, B = 100mM}
sample s4 {1mL, 298.2K} {A = 10mM, B = 10mM}
sample s7 {1mL, 298.2K} {A = 1mM, B = 1mM}
sample s10 {1mL, 298.2K} {A = 100uM, B = 100uM}
```

# Digital Microfluidics Compiler

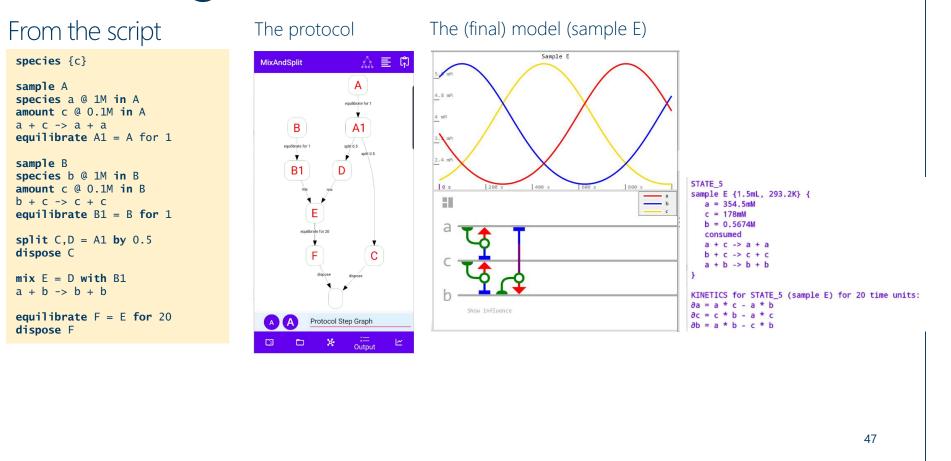
- Mix, split, equilibrate, dispose
- Automatic routing no geometrical information
- Hot/cold zones



### <= Demo: MixAndSplit



### Extracting the Model and the Protocol



#### Extracting the Hybrid Transition System The full story (Hybrid system) From the script species {c} ☆ 🔳 🛱 ☆ ≡ 印 点 🗉 🛱 MixAndSplit ☆ ≡ ₿ MixAndSplit MixAndSplit MixAndSplit TRANSITION [STATE\_1 (equilibrate B1 := B for 1)=> STATE\_2] TRANSITION [STATE\_3 (dispose C)=> STATE\_4] STATE\_2 sample A1 (InL, 293.2K) ( a = 1.064M c = 36.38mM consumed a + c -> a + a A, B STATE\_4 sample B1 {1mL, 293.2K} { b = 0.8512M c = 248.8mM consumed b + c -> c + c sample A STATE\_0 sample A {1mL, 293.2K} { species a @ 1M in A equilibrate A1 := A for 1 a = 1M c = 100mM + amount c @ 0.1M in A b = 0 = 0 = 0 sample B1 {1mL, 293.2K} { b = 0.8512M c = 2.48.8mM consumed b + c => c + c consumed }, sample D (500µL, 293.2K) { a = 1.064M c = 36.38mM consumed a + c -> a + a B, A1 a + c -> a + a a + c -> a + a equilibrate (51 := 8 for 1 sample B {1mL, 293.2K} { equilibrate A1 = A for 1 b = 1M c = 100mM . TRANSITION [STATE\_2 (split C, D := A1 by 0.5)=> STATE\_3] A1, B1 TRANSITION [STATE\_4 (mix E := D, B1)=> STATE\_5] consumed $h + c \rightarrow c + c$ STATE\_3 sample B1 {1mL, 293.2K} { b = 0.8512M c = 248.8MM consumed b + c -> c + c sample B STATE\_5 sample E {1.5mL, 293.2K} { a = 354.5mM b = 0.5574M consumed a + c -> a + a b + c -> c + c a + b -> b + b split C, D :- A1 by 0.5 species b @ 1M in B . KINETICS for STATE\_0 (sample A) for 1 time units: B1, C, D $\partial a = a * c$ $\partial c = -a * c$ amount c @ 0.1M in B }, sample C {500µL, 293.2K} { a = 1.064M c = 36.38mM b + c -> c + cdispose C TRANSITION • equilibrate B1 = B for 1 [STATE\_0 (equilibrate A1 := A for 1)=> STATE\_1] a + c -> a + a KINETICS for STATE\_5 (sample E) for 20 time units: $\partial a = a + c - a + b$ $\partial c = c + b - a + c$ $\partial b = a + b - c + b$ B1, D }, sample D {500µL, 293.2K} { a = 1.064M c = 36.38mM STATE\_1 sample B {1mL, 293.2K} { b = 1M c = 100mM mix E = D. B1 split $C_D = A1$ by 0.5 consumed a + c -> a + a TRANSITION ISTATE 5 (equilibrate F := E for 20)=> STATE 61 . dispose C E consumed A A System Equations A A System Equations b + c -> c + c equilibrate F := E for 20 sample A1 {1mL, 293.2K} { i≡ Outpu mix E = D with B1 + a = 1.064M c = 36.38mM a + b -> b + bF consumed STATE\_6 sample F (1.5mL, 293.2K) { a = 0.5267M c = 167.6mM b = 405.7mM consumed a + c -> a + a equilibrate F = E for 20 × KINETICS for STATE\_1 (sample B) for 1 time units: consumed $a + c \rightarrow a + a$ $b + c \rightarrow c + c$ $a + b \rightarrow b + b$ dispose F $\partial b = -b * c$ $\partial c = b * c$ Protocol State Graph TRANSITION [STATE\_6 (dispose F)=> STATE\_7] A System Equations А STATE\_7 **11** × ~ × Output Output A A System Equations 48

# Kaemika: Extra features

# Extra features

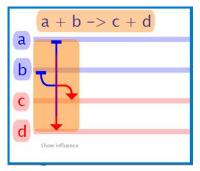
### • General kinetic rates

- Fractions, rational powers, exponentials, trigonometry. E.g., x -> y {{ 1/x }}
- $\cdot$  Work with both deterministic and stochastic simulation and equation-extraction
- Even triggers (discontinuous waveforms)
- Direct ODE notation
  - Instead of a reaction, just write an ODE like  $\partial x = s \cdot y s \cdot x$
  - This is translated to the reaction  $\emptyset$  -> x {{s · y s · x}} using general kinetic rates
- Timeflows (trajectories as first-class values)
  - Programmable plot reports (e.g., var(2 · a 3 · b))
  - · Capture timeflow outputs to combine (e.g., avg) and re-plot/export them later
- Mass action compiler
  - Turn *any* elementary ODE system (with fractions, rational powers, exponentials, trigonometry) into an equivalent system of pure mass action reactions.
- Programmable random numbers and distributions
  - As in MIT's Omega probabilistic language, with rejection sampling.
- Export
  - SBML, ODE, Bitmap, SVG, GraphViz

### Reaction scores (graphical representation of reaction networks)

Horizonal lines: species. Vertical stripes: reactions.

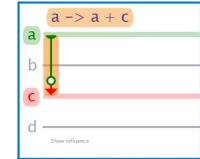
Reactants and products



Repeated species

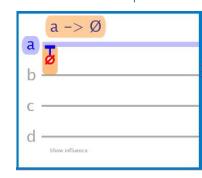
2b -> c + d a b С d Show influence

#### Catalyst but no reactants



#### Reactants but no products

Blue: reagents. Red: products. Green: catalysts.



Catalyst but no products

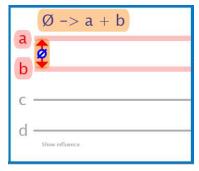
a

b

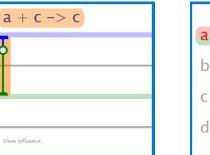
С

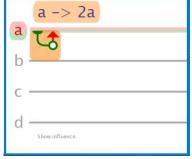
d

Products but no reactants

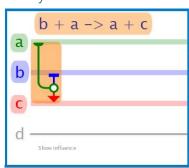


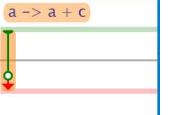
#### Autocatalyst





#### Catalyst

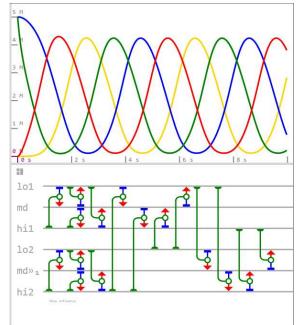


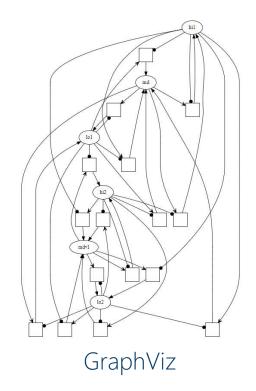


# Reaction Scores vs. Reaction Graphs

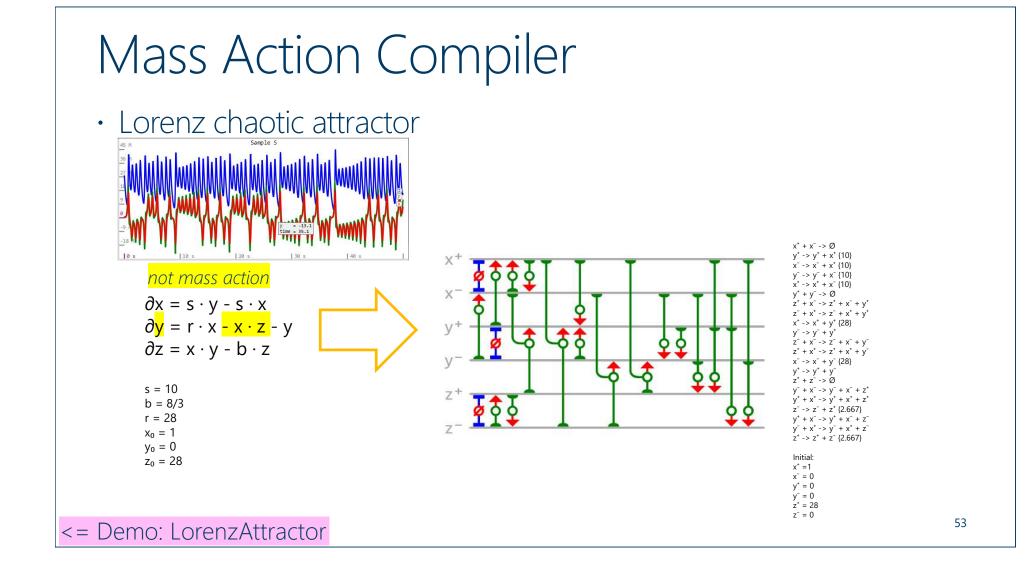
### • 2AM Oscillator

hi1 + md -> 2hi1 lo1 + hi1 -> lo1 + md lo1 + md -> 2lo1 hi2 + lo1 -> hi2 + md {0.5} hi2 + md -> hi2 + hi1 {0.5} lo2 + hi1 -> lo2 + md {0.5} lo2 + md -> lo2 + lo1 {0.5} hi2 + lo2 -> hi2 + md»<sub>1</sub> hi2 + md»<sub>1</sub> -> 2hi2 lo2 + hi2 -> lo2 + md»<sub>1</sub> lo2 + md»<sub>1</sub> -> 2lo2 lo1 + lo2 -> lo1 + md»<sub>1</sub> {0.5} hi1 + md»<sub>1</sub> -> lo1 + hi2 {0.5} hi1 + md»<sub>1</sub> -> hi1 + lo2 {0.5} hi1 + md»<sub>1</sub> -> hi1 + lo2 {0.5}





<= Demo: 2AM Oscillator

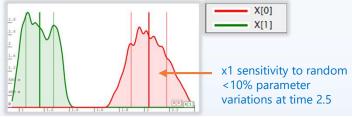


### Global Sensitivity Analysis (of a Lotka-Volterra system)

```
function f(number r1 r2 r3) {
    define
    sample S
    species x1 @ 0.66 M in S
    species x2 @ 0.44 M in S
    x1 -> x1 + x1 {r1}
    x1 + x2 -> x2 + x2 {r2}
    x2 -> Ø {r3}
    equilibrate S for 2.5
    yield [observe(x1,S), observe(x2,S)]
}
```

```
random X(omega w) {
f(1+(w(0)-0.5)/10, 1+(w(1)-0.5)/10, 1+(w(2)-0.5)/10)
```

#### draw 2000 from X



<- A function f to run one simulation (ri are the input parameters to be perturbed)

**Advanced Scripting** 

- <- define D yield E returns the value of E after executing the statements D
- <- Make a new sample S to contain species and reactions for simulation
- <- Lotka-Volterra prey species x1 (initial conditions could be a parameter as well)
- <- Lotka-Volterra predator species x2
- <- Prey reproduces, with perturbed rate r1
- <- Predator eats prey, with perturbed rate r2
- <- Predator dies, with perturbed rate r3
- <- Simulate the system up to time 2.5 (first peak of the oscillation)
- <- Return the output concentrations of x1,x2 from S at time 2.5 as pairs
- <- Create a bivariate random variable X over uniform[0..1) sample spaces w(i)
- <- producing random instances  $f(1+e_1, 1+e_2, 1+e_3) = [x_1,x_2]_{e_1,e_2,e_3,t=2.5}$ with e1, e2, e3 being 10% independent perturbations of the parameters
- <- Produce a density plot of 2000 instances drawn from X i.e. a plot of the distributions of X[0]=x1 and X[1]=x2 at time 2.5 vertical bars are mean and standard deviation

N.B., consider also exporting your Kaemika model to SBML and use the Sobol' method of global sensitivity analysis in e.g. Copasi.

# Conclusions

**Experimental biological protocols with formal semantics** Alessandro Abate, Luca Cardelli, Marta Kwiatkowska, Luca Laurenti, Boyan Yordanov. CMSB 2018.

Kaemika app - Integrating protocols and chemical simulation Luca Cardelli. CMSB 2020.

Kaemika User Manual http://lucacardelli.name/Papers/Kaemika%20User%20Manual.pdf

### Integrated modeling

Of chemical reaction networks and protocols How the Kaemika app supports it Why it needs a *new language* for smooth integration

### Closed-loop modeling, experimentation and analysis

For complete lab automation To "scale up" the scientific method

### Thanks to:

Gold (parser generator) OSLO (ODE simulator) C#/Xamarin (IDE) App store reviewers

### NO thanks to:

XAML (general obfuscator) App store certificates Dark mode support